

Lecture 13: Inductive Program Synthesis

Cognitive Systems - Machine Learning

Part III: Learning Programs and Strategies

Inductive Programming, Functional Programs, Program Schemes, Traces, Folding

last change January 14, 2015

Outlook

- Learning complex rules over symbolic structures
- Program Synthesis
- Inductive Programming: Summers' Thesis
- Igor2
- Application of Igor2 to Cognitive Problems

Knowledge Level Learning

- opposed to low-level (statistical) learning
- learning as generalization of symbol structures (rules) from experience
- “white-box” learning: learned hypotheses are verbalizable, can be inspected, communicated

Examples:

- **Classification/Concepts:**
`IF odor=almond THEN poisonous`
`IF cap-shape=conical & gill-color=grey THEN poisonous`
- **Recursive Concepts:**
`ancestor(X, Y) :- parent(X, Y) .`
`ancestor(X, Y) :- parent(X, Z) , ancestor(Z, Y) .`
- **Simple action rules:**
`IF distance-to-object < threshold THEN STOP`
- **Recursive action rules:**
 e.g., Tower of Hanoi (to be seen later)

Approaches to Symbolical Learning

Machine Learning Approaches

- Grammar Inference
- Decision Tree Algorithms
- Inductive Logic Programming
- Evolutionary Programming
- Inductive (functional) Programming

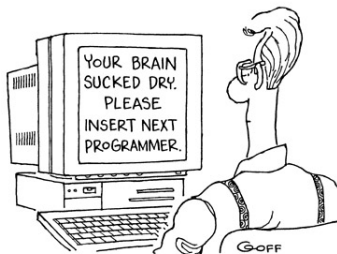
Inductive Programming

- Very special branch of machine learning
- Learning programs from *incomplete* specifications, typically I/O examples or constraints

Program Synthesis

Automagic Programming

- Let the computer program itself
- Automatic code generation from (non-executable) specifications very high level programming
- Not intended for software development in the large but for semi-automated synthesis of functions, modules, program parts



Approaches to Program Synthesis

Deductive and transformational program synthesis

- Complete formal specifications (vertical program synthesis)
- e.g. KIDS (D. Smith)
- High level of formal education is needed to write specifications
- Tedious work to provide the necessary axioms (domain, types, ...)
- Very complex search spaces

$$\forall x \exists y \ p(x) \rightarrow q(x, y)$$

$$\forall x \ p(x) \rightarrow q(x, f(x))$$

Example

*last(l) \Leftarrow find z such that for some y, $l = y \circ [z]$
 where *islist(l)* and $l \neq []$ (Manna & Waldinger)*

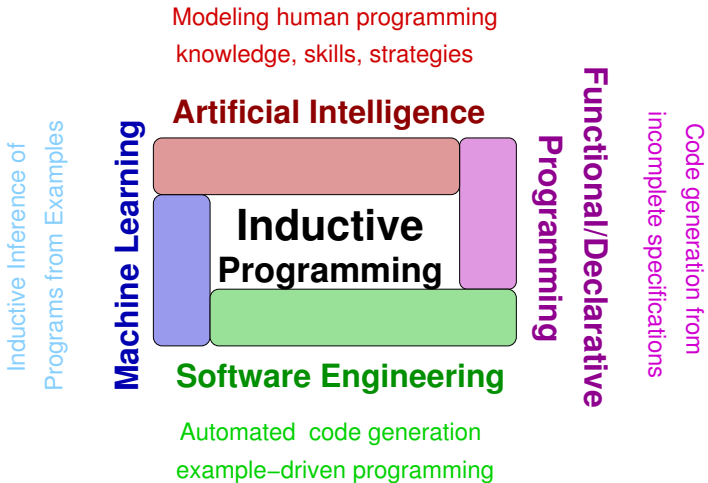
Approaches to Program Synthesis

Inductive program synthesis

- Very special branch of machine learning
- Learning programs from *incomplete* specifications, typically I/O examples or constraints
- Inductive programming (IP) for short

(Flener & Schmid, AI Review, 29(1), 2009; Encyclopedia of Machine Learning, 2010; Schmid, Kitzelmann & Plasmeijr, AAIP 2009)

IP – Contributing Areas



Inductive Programming Example

Learning `last`

I/O Examples

```
last [a]           = a
last [a,b]        = b
last [a,b,c]      = c
last [a,b,c,d]    = d
```

Generalized Program

```
last [x]          = x
last (x:xs)       = last xs
```

Some Syntax

```
-- sugared
[1,2,3,4]
-- normal infix
(1:2:3:4:[])
-- normal prefix
((:) 1
  ((:) 2
    ((:) 3
      ((:) 4
        []))))
```

Inductive Programming – Basics

IP is search in a class of programs (hypothesis space)

Program Class characterized by:

Syntactic building blocks:

- **Primitives**, usually data constructors
- **Background Knowledge**, additional, problem specific, user defined functions
- **Additional Functions**, automatically generated

Restriction Bias

syntactic restrictions of programs in a given language

Result influenced by:

Preference Bias

choice between syntactically different hypotheses

Inductive Programming – Approaches

- Typical for declarative languages (LISP, PROLOG, ML, HASKELL)
- Goal: finding a program which covers *all* input/output examples *correctly* (no PAC learning) and (recursively) generalizes over them
- Two main approaches:
 - **Analytical, data-driven:**
detect regularities in the I/O examples (or traces generated from them) and generalize over them (folding)
 - **Generate-and-test:**
generate syntactically correct (partial) programs, examples only used for testing

Inductive Programming – Approaches

Generate-and-test approaches

- ILP (90ies): FFOIL (Quinlan) (sequential covering)
- evolutionary: ADATE (Olsson)
- enumerative: MAGICHASKELLER (Katayama)
- also in functional/generic programming context: automated generation of instances for data types in the model-based test tool G \forall st (Koopmann & Plasmeijer)

Inductive Programming – Approaches

Analytical Approaches

- Classical work (70ies–80ies):
THESYS (Summers), Biermann, Kodratoff
learn linear recursive Lisp programs from traces
- ILP (90ies):
Golem, Progol (Muggleton), Dialogs (Flener)
inverse resolution, Θ -subsumption, schema-guided
- IGOR1 (Schmid, Kitzelmann; extension of THESYS)
IGOR2 (Kitzelmann, Hofmann, Schmid)

Summers' Thesys

Summers (1977), A methodology for LISP program construction from examples, Journal ACM

Two Step Approach

- Step 1: Generate traces from I/O examples
- Step 2: Fold traces into recursion

Generate Traces

- Restriction of input and output to nested lists
- Background Knowledge:
 - Partial order over lists
 - Primitives: `atom`, `cons`, `car`, `cdr`, `nil`
- Rewriting algorithm with unique result for each I/O pair: characterize I by its structure (lhs), represent O by expression over I (rhs)

↔ restriction of synthesis to structural problems over lists (abstraction over elements of a list) not possible to induce `member` or `sort`

Example: Rewrite to Traces

I/O Examples

$\text{nil} \rightarrow \text{nil}$

$(A) \rightarrow ((A))$

$(A B) \rightarrow ((A) (B))$

$(A B C) \rightarrow ((A) (B) (C))$

Traces

$F_L(x) \leftarrow$

$$\begin{aligned}
 & (\text{atom}(x) \rightarrow \text{nil}, \\
 & \text{atom}(\text{cdr}(x)) \rightarrow \text{cons}(x, \text{nil}), \\
 & \text{atom}(\text{caddr}(x)) \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cdr}(x), \text{nil})), \\
 & \text{T} \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cons}(\text{caddr}(x), \text{nil}), \\
 & \quad \text{cons}(\text{caddr}(x), \text{nil})))
 \end{aligned}$$

Example: Deriving Fragments

Unique Expressions for Fragment (A B)

$(x, (A\ B)),$
 $(\text{car}[x], A),$
 $(\text{cdr}[x], (B)),$
 $(\text{cadr}[x], B),$
 $(\text{caddr}[x], ())$

Combining Expressions

$((A)\ (B)) = \text{cons}[(A); ((B))] = \text{cons}[\text{cons}[A, ()]; \text{cons}[(B), ()]].$

Replacing Values by Functions

$\text{cons}[\text{cons}(\text{car}[x]; ()); \text{cons}[\text{cdr}[x]; ()]]$

Folding of Traces

- Based on a program scheme for linear recursion (restriction bias)
- Synthesis theorem as justification
- Idea: inverse of fixpoint theorem for linear recursion
- Traces are k th unfolding of an unknown program following the program scheme
- Identify differences, detect recurrence

$$\begin{aligned}
 F(x) \leftarrow & (p_1(x) \rightarrow f_1(x), \\
 & \dots, \\
 & p_k(x) \rightarrow f_k(x), \\
 & T \rightarrow C(F(b(x)), x))
 \end{aligned}$$

Example: Fold Traces

kth unfolding

$$\begin{aligned}
 F_L(x) \leftarrow & \text{ (atom}(x) \rightarrow \text{nil,} \\
 & \text{atom}(\text{cdr}(x)) \rightarrow \text{cons}(x, \text{nil}), \\
 & \text{atom}(\text{caddr}(x)) \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cdr}(x), \text{nil})), \\
 & \text{T} \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cons}(\text{caddr}(x), \text{nil}), \\
 & \quad \text{cons}(\text{caddr}(x), \text{nil}))))
 \end{aligned}$$

Differences:

$$p_2(x) = p_1(\text{cdr}(x))$$

$$p_3(x) = p_2(\text{cdr}(x))$$

$$p_4(x) = p_3(\text{cdr}(x))$$

$$f_2(x) = \text{cons}(x, f_1(x))$$

$$f_3(x) =$$

$$\text{cons}(\text{cons}(\text{car}(x), \text{nil}), f_2(\text{cdr}(x)))$$

$$f_4(x) =$$

$$\text{cons}(\text{cons}(\text{car}(x), \text{nil}), f_3(\text{cdr}(x)))$$

Recurrence Relations:

$$p_1(x) = \text{atom}(x)$$

$$p_{k+1}(x) = p_k(\text{cdr}(x)) \text{ for } k = 1, 2, 3$$

$$f_1(x) = \text{nil}$$

$$f_2(x) = \text{cons}(x, f_1(x))$$

$$f_{k+1}(x) = \text{cons}(\text{cons}(\text{car}(x), \text{nil}), f_k(\text{cdr}(x)))$$

$$\text{for } k = 2, 3$$

Example: Fold Traces

kth unfolding

$$\begin{aligned}
 F_L(x) \leftarrow & \text{ (atom}(x) \rightarrow \text{nil,} \\
 & \text{atom}(\text{cdr}(x)) \rightarrow \text{cons}(x, \text{nil}), \\
 & \text{atom}(\text{cddr}(x)) \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cdr}(x), \text{nil})), \\
 & \text{T} \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{cons}(\text{cons}(\text{cadr}(x), \text{nil}), \\
 & \quad \text{cons}(\text{cddr}(x), \text{nil}))))
 \end{aligned}$$

Folded Program

$$\begin{aligned}
 \text{unpack}(x) \leftarrow & \text{ (atom}(x) \rightarrow \text{nil,} \\
 & \text{T} \rightarrow \text{u}(x)) \\
 \text{u}(x) \leftarrow & \text{ (atom}(\text{cdr}(x)) \rightarrow \text{cons}(x, \text{nil}), \\
 & \text{T} \rightarrow \text{cons}(\text{cons}(\text{car}(x), \text{nil}), \text{u}(\text{cdr}(x))))
 \end{aligned}$$

Summers' Synthesis Theorem

- Based on fixpoint theory of functional program language semantics.
(Kleene sequence of function approximations: a partial order can be defined over the approximations, there exists a supremum, i.e. least fixpoint)
- Idea: If we assume that a given trace is the k -th unfolding of an unknown linear recursive function, then there must be regular differences which constitute the stepwise unfoldings and in consequence, the trace can be generalized (folded) into a recursive function

Illustration of Kleene Sequence

Defined for no input

$$U^0 \leftarrow \Omega$$

Defined for empty list

$$U^1 \leftarrow \begin{array}{l} (atom(x) \rightarrow nil, \\ T \rightarrow \Omega) \end{array}$$

Defined for empty list and lists with one element

$$U^2 \leftarrow \begin{array}{l} (atom(x) \rightarrow nil, \\ atom(cdr(x)) \rightarrow cons(x, nil), \\ T \rightarrow \Omega) \end{array}$$

... Defined for lists up to n elements

The IP System IGOR2

Inductive

- induces programs from I/O examples
- inspired by Summers' THESYS system
- successor of IGOR1 (Schmid, 2001)

Analytical

- data-driven
- finds recursive generalization by analyzing I/O examples
- integrates best first search

Functional

- learns functional programs
- first prototype in MAUDE by Emanuel Kitzelmann
- in HASKELL and extended (general *fold*) by M. Hofmann

Example

reverse

I/O Example

```
reverse [] = []           reverse [a,b] = [b,a]
reverse [a] = [a]        reverse [a,b,c] = [c,b,a]
```

Generalized Program

```
reverse [] = []
reverse (x:xs) = last (x:xs) : reverse (init (x:xs))
```

Automatically induced functions (*renamed* from $f1$, $f2$)

```
last [x] = x           init [a] = []
last (x:xs) = last xs  init (x:xs) = x:(init xs)
```

Input

Data-type Definitions

```
data [a] = [] | a:[a]
```

Target Function

```
reverse :: [a] -> [a]
reverse [] = []
reverse [a] = [a]
reverse [a,b] = [b,a]
reverse [a,b,c] = [c,b,a]
```

Background Knowledge

```
snoc :: [a] -> a -> [a]
snoc [] x = [x]
snoc [x] y = [x,y]
snoc [x,y] z = [x,y,z]
```

- Input must be the first k I/O examples (wrt to input data type)
- Background knowledge is *optional*

Output

Set of (recursive) equations which cover the examples

`reverse` Solution

`reverse [] = []`

`reverse (x:xs) = snoc (reverse xs) x`

Restriction Bias

- Subset of HASKELL or MAUDE
- Case distinction by *pattern matching*
- Syntactical restriction: patterns are not allowed to unify

Preference Bias

- Minimal number of case distinctions

Basic Idea

- Search a rule which explains/covers a (sub-) set of examples
- Initial hypothesis is a single rule which is the least general generalization (anti-unification) over all examples

Example Equations

reverse [a] = [a]

reverse [a,b] = [b,a]

Initial Hypothesis

reverse (x:xs) = (y:ys)

Hypothesis contains *unbound* variables in the body!

Basic Idea cont.

Initial Hypothesis

$$\text{reverse } (x:xs) = (y:ys)$$

Unbound variables are cue for induction.

Three Induction Operators (to apply simultaneously)

- 1 **Partitioning** of examples
 \rightsquigarrow *Sets* of equations divided by case distinction
- 2 Replace right-hand side by **program call** (recursive or background)
- 3 Replace sub-terms with unbound variables by to be induced **sub-functions**

Kitzelmann & Schmid, JMLR, 7, 2006; Kitzelmann, LOPSTR, 2008; Kitzelmann doctoral thesis 2010

Some Empirical Results (Hofmann et al. AGI'09)

	<i>isort</i>	<i>reverse</i>	<i>weave</i>	<i>shiftr</i>	<i>mult/add</i>	<i>allodds</i>
ADATE	70.0	78.0	80.0	18.81	—	214.87
FLIP	×	—	134.24 [⊥]	448.55 [⊥]	×	×
FFOIL	×	—	0.4 [⊥]	< 0.1 [⊥]	8.1 [⊥]	0.1 [⊥]
GOLEM	0.714	—	0.66 [⊥]	0.298	—	0.016 [⊥]
IGOR II	0.105	0.103	0.200	0.127	⊙	⊙
MAGH.	0.01	0.08	⊙	157.32	—	×

	<i>lasts</i>	<i>last</i>	<i>member</i>	<i>odd/even</i>	<i>multlast</i>
ADATE	822.0	0.2	2.0	—	4.3
FLIP	×	0.020	17.868	0.130	448.90 [⊥]
FFOIL	0.7 [⊥]	0.1	0.1 [⊥]	< 0.1 [⊥]	< 0.1
GOLEM	1.062	< 0.001	0.033	—	< 0.001
IGOR II	5.695	0.007	0.152	0.019	0.023
MAGH.	19.43	0.01	⊙	—	0.30

— not tested × stack overflow ⊙ timeout ⊥ wrong

Application of IGOR2 to Cognitive Problems

(Schmid & Kitzelmann, CSR, 2011)

Problem Solving

- Clearblock (4 examples, 0.036 sec)
- Rocket (3 examples, 0.012 sec)
- Tower of Hanoi (3 examples, 0.076 sec)
- Car Park (4 examples, 0.024 sec)
- Blocks-world Tower (9 examples, 1.2 sec)

Recursive Concepts

- Ancestor (9 examples, 10.1 sec)

Syntactic Rules

- Phrase structure grammar (3 examples, 0.072 sec)

$$S \rightarrow NP VP \quad NP \rightarrow d n \quad VP \rightarrow v NP \mid v S$$

Learning Tower of Hanoi

Input to IGOR2

```

eq Hanoi(0, Src, Aux, Dst, S) =
  move(0, Src, Dst, S) .
eq Hanoi(s 0, Src, Aux, Dst, S) =
  move(0, Aux, Dst,
    move(s 0, Src, Dst,
      move(0, Src, Aux, S))) .
eq Hanoi(s s 0, Src, Aux, Dst, S) =
  move(0, Src, Dst,
    move(s 0, Aux, Dst,
      move(0, Aux, Src,
        move(s s 0, Src, Dst,
          move(0, Dst, Aux,
            move(s 0, Src, Aux,
              move(0, Src, Dst, S))))))) .

```

Induced Tower of Hanoi Rules (3 examples, 0.076 sec)

```

Hanoi(0, Src, Aux, Dst, S) = move(0, Src, Dst, S)
Hanoi(s D, Src, Aux, Dst, S) =
  Hanoi(D, Aux, Src, Dst,
    move(s D, Src, Dst,
      Hanoi(D, Src, Dst, Aux, S)))

```

Applying IGOR2 to Number Series Induction

- IGOR2 is designed as IP system
- Generalization over traces/streams of observations to productive rules
- Learning from few, small examples, generalization to n
- IGOR2 as a “cognitive rule acquisition device” (ct. Chomsky’s LAD)?
- Detterman challenge 2011: An AI system should be able to solve a variety of different problems without being engineered to each special application
- IGOR2 can solve different cognitive problems “from the shelf”!
- Further example: IQ test problems – number series (J. Hofmann, Kitzelmann, Schmid, KI’14)

Applicability of IGOR2 to number series problems

- Crucial: How to represent number series problems as input for IGOR2

(1) Input List – Output Successor Value

eq Plustwo((s 0) nil) = s^3 0

eq Plustwo((s^3 0) (s 0) nil) = s^5 0

eq Plustwo((s^5 0) (s^3 0) (s 0) nil) = s^7 0

(2) Input Position – Output List

eq Plustwo(s 0) = (s 0) nil

eq Plustwo(s^2 0) = (s^3 0) (s 0) nil

eq Plustwo(s^3 0) = (s^5 0) (s^3 0) (s 0) nil

eq Plustwo(s^4 0) = (s^7 0) (s^5 0) (s^3 0) (s 0) nil

(3) Input Position – Output Value

eq Plustwo(s 0) = s 0

eq Plustwo(s s 0) = s s s 0

eq Plustwo(s s s 0) = s s s s s 0

eq Plustwo(s s s s 0) = s s s s s s s 0

Series examples

1 2 3 12 13 14 23	$f(n-3) + 11$	+, 1, small, large, linear, $n-3$, const
1 2 3 5 8	$f(n-1) + f(n-2)$	+, 1, small, small, comp, $n-1/n-2$, const
6 7 8 18 21 24 54	$f(n-3) \times 3$	\times , 1, small, small, linear, $n-3$, const
3 4 12 48 576	$f(n-1) \times f(n-2)$	\times , 1, small, sm/lrg, comp, $n-1/n-2$, const
5 10 30 120 600	$f(n-1) \times n$	\times , 1, large, large, linear, $n-1$, pos
15 15 16 15 15 16 15	$f(n-3)$	=, 1, large, large, linear, $n-3$, const

Combining IP, Planning, Analogy

- Learning from planning: generate action traces with an AI planning system and learn a recursive rule set
- Use already known recursive rule sets to solve structurally similar problems (analogy)
- Generalization over base and target problems results in a hierarchy of program/problem solving schemes

Summary

- Learning can be high-level, on symbolic representations or low-level (e.g., feature vectors)
- Reinforcement learning addresses learning of control strategies/policies based on a statistical approach
- Inductive programming addresses learning of recursive rule sets from few examples or from traces
- IP is part of automated programming research, it complements deductive approaches
- The first IP system was Summers's THESYS, Summers also provided the formal foundation (Synthesis Theorem)
- IGOR2 is a current approach to IP, it is more general than THESYS and incorporates aspects from ILP (use of background knowledge, function invention)
- IP approaches are often general enough to be applied to other strategy learning domains, e.g., from cognition

Summary

- In contrast to classical ML, IP needs only few and only positive examples
- IP is nevertheless *supervised* (learning functions: negative examples are present implicitly)

Learning Terminology

Thesys, Igor2

Supervised Learning	unsupervised learning
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Approaches:

Concept / Classification	Policy Learning
symbolic	statistical / neuronal network
inductive	analytical

Learning Strategy: Data: Target Values:	learning from examples input/output examples or traces Recursive rule sets
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